

Homework 4

Due: March 1, 2023

All homeworks are due at 11:59 PM on Gradescope.

Please do not include any identifying information about yourself in the handin, including your Banner ID.

Be sure to fully explain your reasoning and show all work for full credit.

Problem 1

Recall that an *equivalence relation* is one that is reflexive, symmetric, and transitive.

Determine whether or not each of the following relations is an equivalence relation. Be sure to justify your answers.¹ In particular: if a relation is *not* an equivalence relation, which of the above properties does it not satisfy?

- a. The relation R on \mathbb{Z} defined by the set of ordered pairs of integers:

$$R = \{(a, b) \mid |a - b| \geq 3\}.$$

- b. The relation S on \mathbb{R}^2 defined by the set of ordered pairs of coordinates:

$$S = \{(a, b) \mid \|a\| = \|b\|\},$$

where $\|a\|$ is the distance from a to the origin in \mathbb{R}^2 (\mathbb{R}^2 is the set of ordered pairs (x, y) where $x, y \in \mathbb{R}$, also known as the set of points in the plane, and the distance from a point (x, y) to the origin is defined as $\sqrt{x^2 + y^2}$.)

In addition to your proof, answer the following: given a fixed point $p \in \mathbb{R}^2$, the collection of all points related to p gives what familiar geometric object? That is, what is $\{x \mid x S p\}$?

- c. Let $A = \{a, b, c, d\}$. Let T be the relation on A with the graph:

$$\{(a, a), (b, b), (c, c), (a, b), (b, c), (a, c), (b, a), (c, b), (c, a)\}.$$

¹You might find this resource [here](#) helpful for this question.

Problem 2

The CS22 plant nursery is now open and quickly growing in popularity! The first customer came in and bought 1 plant. The second bought 8. The third bought 27. The pattern has continued: the k th customer bought k^3 plants.

This is a lot of plants, and to keep our inventory up, we want to predict how many plants in total we will have sold after the n th customer has made their purchase. Call this number of total sales $S(n)$.

Our staff has made some guesses for the value of $S(n)$:

1. $S(n) = n^3 + (n - 1)^3$

2. $S(n) = \frac{1}{4}n^2(n + 1)^2$

3. $S(n) = 8^{n-1}$

- a. Assuming that one of the guesses is correct, which one do you believe? Briefly justify your answer. (You do not need to prove it; that's the next part.)
- b. Prove, using induction, that your chosen formula is correct. You should follow the "recipe" for an induction proof, stating the induction predicate clearly.

Problem 3

This problem is a Lean question!

This homework question can be found by navigating to `BrownCs22/Homework/Hw4.lean` in the directory browser on the left of your screen in Gitpod. The comment at the top of that file provides more detailed instructions.

NOTE: the submission method is different this week! Now, you can submit a .pdf file, a .tex file, and a .lean file all to the same assignment. The autograder will check your .lean file as normal, and also make sure you have submitted the .tex source for the rest of your solutions.

Question 3 Autograded on Lean/Gradescope.



Problem 4 (Mind Bender — *Extra Credit*)

Mind Benders are extra credit problems intended to be more challenging than usual homework problems and are an exploration into a topic not covered in lecture. This week, we'll think about induction over sets other than \mathbb{N} .

First of all: what is \mathbb{N} , exactly? Our definition $\{0, 1, 2, \dots\}$ wasn't very precise. Here's a more careful attempt. We'll say a set S is *number-like* if it satisfies the following properties:

1. $0 \in S$.
2. For every $n \in S$, $\text{successor}(n) \in S$.
3. The **successor** function is injective.
4. For every $n \in S$, $\text{successor}(n) \neq 0$.
5. Every element of S is either 0 or the successor of some other element of S . (In other words, rules 1 and 2 are the only ways to “create” natural numbers. In yet more words, the range of **successor** is $S \setminus \{0\}$.)

We then define \mathbb{N} to be the *smallest* number-like set: that is, the intersection of all number-like sets. Note that 0 is in this intersection, since it must be in every number-like set. Similarly for $\text{successor}(0)$, ...

- a. Using this characterization of \mathbb{N} , explain why the principle of induction works to prove facts about all natural numbers. (Note: $n + 1$ is another way to write $\text{successor}(n)$.)
- b. We need each of properties 3-5 to make sure that \mathbb{N} has the shape we want.

Example

Suppose we defined $\mathcal{N} = \{0, 1, 2\}$ with

$$\text{successor}(0) = 1$$

$$\text{successor}(1) = 2$$

$$\text{successor}(2) = 0$$

\mathcal{N} satisfies properties 1, 2, 3, and 5, but fails 4.

Give an example like this of a set and successor function that satisfies every property *except* property 3, and another example that satisfies every property *except* 5.

We can give a similar characterization of *lists* of elements of some set T .

1. The empty list $[]$ is a list of elements of T .
2. For any $t \in T$ and list L of elements of T , $\mathbf{append}(t, L)$ is a list of elements of T : that is, it's the list we get when we *append* t at the beginning of L .
- 3-5. Similar to 3-5 above for \mathbb{N} .

This may not be a definition of lists that you've seen before. We call it, maybe tellingly, an *inductive* (or *algebraic*) definition.

- c. "Translate" properties 3-5 from the \mathbb{N} case so that they make sense in the list case.
- d. Lists can be a kind of mathematical object just like numbers: we may want to prove a theorem about all lists of elements of some set T . Analogous to the principle of induction for \mathbb{N} , there's a principle of induction for lists. State this principle of induction.
- e. We say a *subsequence* is a non-necessarily contiguous subset of a list L . For example, if L were the list $[1, 2, 3, 4, 5]$, then one such subsequence could be $[2, 4, 5]$. Prove, by inducting on the structure of lists, that the number of subsequences of a list of length n (containing all distinct elements) is 2^n .
- f. Can you think of any other sets/structures that can be defined "inductively" like this?

Note: these sets show up all the time in certain areas of computer science. Some programming languages *only* allow you to define new datatypes as inductive sets² (or "inductive types")! What if you write a complicated program that takes inputs from one of these datatypes, and want to prove that your program has certain behavior? You guessed it: time for induction.

²For those who have taken a computer science course that has introduced functional programming, \mathbf{append} may look very similar to \mathbf{link} or \mathbf{cons} !